## Example:

$A+B \quad(A=1011)$ and $(B=0011)$

| Subscript i | 3 | 2 | 1 | 0 |  | $C_{0}=0$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Input Carry | 0 | 1 | 1 | 0 | $C_{i}$ |  |
| A | 1 | 0 | 1 | 1 |  |  |
| + | 1 | 0 | 1 | 1 | $A_{i}$ |  |
| B | 0 | 0 | 1 | 1 | $\boldsymbol{B}_{\boldsymbol{i}}$ |  |
| Sum | 1 | 1 | 1 | 0 | $S_{i}$ |  |
| Output Carry | 0 | 0 | 1 | 1 | $C_{i+1}$ |  |

## Carry Propagation

$>$ The addition of $A+B$ binary numbers in parallel implies that all the bits of $A$ and $B$ are available for computation at the same time.
$>$ As in any combinational circuit, the signal must propagate through the gates before the correct output sum is available.
> The output will not be correct unless the signals are given enough time to propagate through the gates connected form the input to the output.
$>$ The longest propagation delay time in an adder is the time it takes the carry to propagate through the full adders.


Full Adder with P and G
$>$ The signal form the carry input $C_{i}$ to the output carry $C_{i+1}$ propagates through an $A N D$ gate and an $O R$ gate, which equals 2 gate levels.

- If there are 4 full adders in the binary adder, the output carry $C_{4}$ would have $2 \times 4=8$ gate levels, form $C_{0}$ to $C_{4}$
- For an $n$-bit adder, $2 n$ gate levels for the carry to propagate form input to output are required.

